

Languages & Grammars

Syntax vs. Semantics

↑ ↑
Structure meaning.

```
if (5 > 10) {  
    x = x;  
    x = x + 1 - 1;  
    x = x + 0;  
} else {  
    x = x;  
}  
}
```

```
if (x == 0) {  
    ~  
    —  
}  
int *p;  
int i;  
p = i;
```

Example of a formal language

Sentence = noun phrase + verb phrase

noun phrase = article + adjective + noun
or
article + noun

verb phrase = ~~adverb~~ + verb + ~~adverb~~
or
verb

article = "the" or "a"

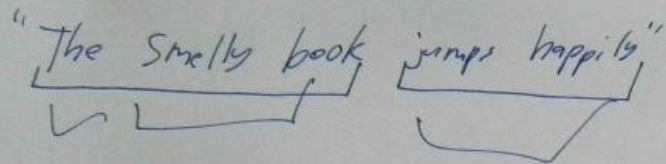
adjective = "large" or "small" or "tall"

adverb = "quickly" or "happily"

noun = "car" or "rabbit" or "textbook"

verb = "runs" or "jumps" or "vanishes"

"The Smelly book jumps happily"



the cat vanishes.

Specifies a language

- make a list of all words.
- list rules for what is valid
- Describe a phrase structure grammar

Definitions

A vocabulary V is a finite set of symbols, $\neq \emptyset$

A word (or, sometimes, a ~~set~~ sequence) over V is a finite-length string of symbols from V .

λ = empty string (word of length 0)

$V^* = \{\text{set of all words over } V\}$

A language over V is a subset of V^* .

A phrase structure grammar $G = (V, T, S, P)$

V = a vocabulary

$T = \{\text{terminal symbols}\} \subseteq V$

S = start symbol $\in V$

$P = \{\text{production rules}\}$

rules for replacing one word with another

EXAMPLES

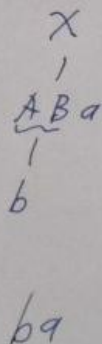
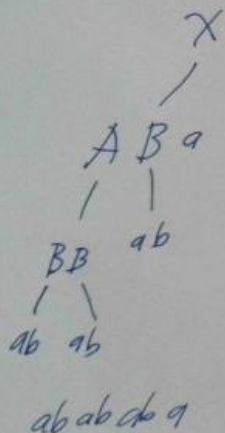
$$G = (V, T, S, P)$$

$$V = \{a, b, A, B, x\}$$

$$T = \{a, b\}$$

$$S = x$$

$$P = \{x \rightarrow ABa, A \rightarrow BB, B \rightarrow ab, AB \rightarrow b\}$$



$$G = (V, T, S, P)$$

$$V = \{s\} \quad T = \{0, 1\} \quad S = s$$

$$P = \{s \rightarrow 11s, s \rightarrow 0\}$$

s
1
0

s
1
11s
1
0

s
1
11s
1
11s
1
0

0

110

11110

1111110

11111110 ...

$$\{(11)^n 0, n \geq 0\}$$

$$V = \{0, 1, A, B, s\} \quad T = \{0, 1\} \quad S = s$$

$$P = \{S \rightarrow 0A, S \rightarrow 1A, A \rightarrow 0B, B \rightarrow 1A, B \rightarrow 1\}$$

$$\begin{array}{c} S \\ | \\ 0A \\ \hline \end{array}$$

$$\begin{array}{c} A \\ | \\ 0B \\ | \\ 1 \\ | \end{array}$$

$$\begin{array}{c} A \\ | \\ 0B \\ | \\ 1A \\ \hline \end{array}$$

$$01$$

$$01A$$

$$0101A$$

$$010101$$

$$010101A$$

$$(01)^n$$

$$n \geq 1$$

$$\{0(01)^n, 1(01)^n \mid n \geq 1\}$$

$$V = \{0, 1, A, B, s\} \quad T = \{0, 1\} \quad S = s$$

$$P = \{S \rightarrow 0A, S \rightarrow 1A, A \rightarrow 0B, B \rightarrow 1A, B \rightarrow 1\}$$

$$\begin{array}{c} S \\ | \\ 0A \\ \hline \end{array}$$

$$\begin{array}{c} A \\ | \\ 0B \\ | \\ 1 \\ | \end{array}$$

$$\begin{array}{c} A \\ | \\ 0B \\ | \\ 1A \\ \hline \end{array}$$

$$01$$

$$01A$$

$$0101A$$

$$010101$$

$$010101A$$

$$(01)^n$$

$$n \geq 1$$

$$\{0(01)^n, 1(01)^n \mid n \geq 1\}$$

$$0010101$$

Constructing a grammar

Suppose $L(G) = \{\lambda, 01, 0011, 000111, 00001111, \dots\}$

$$V = \{0, 1, S\}$$

$$T = \{0, 1\}$$

$$S \vdash S$$

$$P = \{S \rightarrow \lambda, (S \rightarrow 01) (\cancel{S \rightarrow 0011}, \cancel{S \rightarrow 000111}) (S \rightarrow 0S1)\}$$

S
|
0S1
|
λ
01

S
|
0S1
|
0S1
|
01
000111

S
|
0S1
|
01
0011

$$L(G) = \{0^m, m \geq 0\}$$

$$V = \{0, s\} \quad T = \{0\} \quad S = s$$

$$P = \{s \rightarrow 0, s \rightarrow \lambda, s \rightarrow 0s\}$$

$$L(G) = \{0^n, n \geq 1\}$$

$$S \rightarrow 0J$$

$$J \rightarrow \lambda, J \rightarrow 0J$$

$$L(G) = 1^n, n \geq 1$$

$$S \rightarrow 1K$$

$$K \rightarrow \lambda, K \rightarrow \cancel{1}K$$

$$L(G) = \{0^m 1^n, mn \geq 1\}$$

$$V = \{S, 0, 1\}$$

$$T = \{0, 1\}$$

$$S \rightarrow S$$

$$P = \{S \rightarrow AB, A \rightarrow 0J, J \rightarrow \lambda, J \rightarrow 0J, \\ B \rightarrow 1K, K \rightarrow \lambda, K \rightarrow 1K\}$$

$$P = \{S \rightarrow 0S, S \rightarrow 1A, A \rightarrow 1A, A \rightarrow \lambda\} \quad ?$$

Let $w_0 = L Z_0 R$

(w_0, w_1, L, R, Z_0, Z_1
are strings)

Let $w_1 = L Z_1 R$

$G(V, T, S, P)$

If $Z_0 \rightarrow Z_1 \in P$ then

we say "~~that~~ w_1 is directly derivable from w_0 "

$$w_0 \Rightarrow w_1$$

If $w_0 \Rightarrow w_1 \Rightarrow w_2 \Rightarrow w_3 \Rightarrow \dots \Rightarrow w_n$ then

we say " w_n is derivable from w_0 "

$$w_0 \xRightarrow{*} w_n$$

$L(G)$ is "the language generated by G " = {string of terminal symbols
derivable from S }

$$L(G) = \{w \in T^* : S \xRightarrow{*} w\}$$

BNF (Backus-Naur form) for an assignment statement

$AS := \text{var } "=" \text{ exp}$

$\text{var} := \text{letter idstring} / \text{letter}$

$\text{idstring} := \text{letter idstring} / \text{digit idstring} / \text{letter} / \text{digit}$

$\text{letter} := 'A' / 'B' / 'C' / 'D' / \dots / 'Z'$

$\text{digit} := '0' / '1' / '2' / \dots / '9'$

$\text{exp} := \text{~~var~~ exp op exp} / \text{var} / \text{number} / "(" \text{exp} ")"$

$\text{op} := "+" / "-" / "/" / "*" / "\times"$

The syntax of C in Backus-Naur Form

`<translation-unit> ::= {<external-declaration>}*`

`<external-declaration> ::= <function-definition>
| <declaration>`

`<function-definition> ::= {<declaration-specifier>}* <declarator> {<declaration>}* <compound-statement>`

`<declaration-specifier> ::= <storage-class-specifier>
| <type-specifier>
| <type-qualifier>`

`<storage-class-specifier> ::= auto
| register
| static
| extern
| typedef`

`<type-specifier> ::= void
| char
| short
| int
| long
| float
| double
| signed
| unsigned
| <struct-or-union-specifier>
| <enum-specifier>
| <typedef-name>`

`<struct-or-union-specifier> ::= <struct-or-union> <identifier> { {<struct-declaration>}+ }`