

12

$$\binom{12}{3} \binom{9}{4} \binom{5}{5} = \cancel{37} > ? ? 2.$$

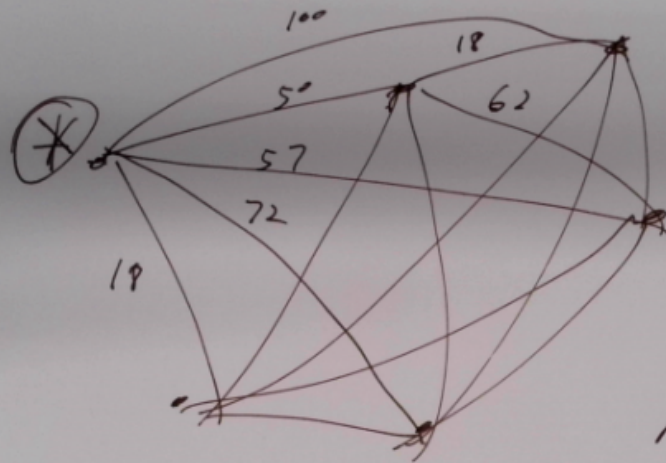
$$\binom{12}{5} \binom{7}{3} \binom{4}{4}$$

$$\binom{12}{5} \binom{7}{4} \binom{3}{3}$$

$$\binom{12}{4} \binom{8}{3} \binom{5}{5}$$

⋮

Traveling Salesperson (TS)



$$\sim h!$$

$$\sim h^n$$

- h
- $n-1$
- $h-2$
- \vdots
- \vdots

- n
- n
 - n
 - n
 - n
 - n
 - \vdots

Complexity

$$n \sim n^2$$

$$\underline{O(n^2)}$$

$O(n^3)$ polynomial time

$O(n)$

$O(n^n)$ TS: $\sim O(n^n)$

Is there a poly-time alg. for TS?

Language & Grammars

Syntax vs. Semantics

```
IF (1 > 2) {  
    x = x;  
    x = x;  
    x = x + 0;  
} else {  
    x = x + 0;  
}
```

The sleeping colorers
green gas is landing.

Sentence = noun phrase + verb phrase

Noun phrase = article + adj + noun
or article + noun

Verb phrase = verb + adverb
or verb

article = "a" or "the"

adj = "large" or "human" or "smells"

noun = "rabbit" or "teacher"

verb = "eats" or "jumps" or "flies"

adverb = "quickly" or "happily"

A Vocabulary V is a finite, non-empty set of symbols.

A word or sentence over V is a finite-length string of elements of V .

λ = empty string.

$V^* = \{\text{all words over } V\}$

A language over V is a subset of V^* .

Q: How to specify a language?

- 1) use rules
- 2) make a full list of all words
- *3) Describe a grammar

A phrase structure grammar $G = (V, T, S, P)$

V : Vocabulary

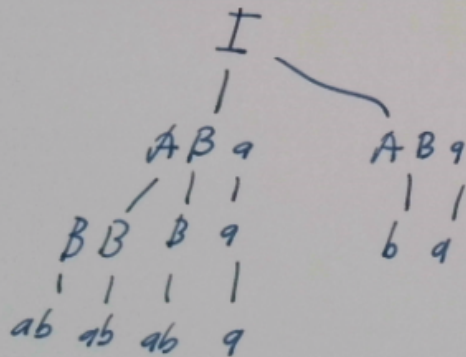
T : Set of terminal symbols $\subseteq V$

S : Start symbol ($S \in V$)

P : Set of production rules

Example: $V = \{a, b, A, B, I\}$ $T = \{a, b\}$ $S = I$

$P = \{ I \rightarrow ABa, A \rightarrow BB, B \rightarrow ab, AB \rightarrow b \}$



abababab

ba

DEFS: Let $w_0 = L z_0 R$

$w_1 = L z_1 R$

$G = (V, T, S, P)$

Suppose $z_0 \rightarrow z_1 \in P$

Then we say " w_1 is directly derivable from w_0 "

$w_0 \Rightarrow w_1$

If $w_0 \Rightarrow w_1 \Rightarrow w_2 \Rightarrow \dots \Rightarrow w_n$

then we say " w_n is derivable from w_0 "

$w_0 \xRightarrow{*} w_n$

DEF: The language generated by G $L(G) =$

$\{ \text{strings of terminals which are derivable from } S \}$

$L(G) = \{ w \in T^* : S \xRightarrow{*} w \}$

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$$V = \{I, 0, 1\}, T = \{0, 1\}, S = I, P = \{I \rightarrow 11I, I \rightarrow 0\} \quad G = (V, T, S, P)$$

$$L(G) = \{0, 110, 11110, 1111110, \dots\} = 1^n 0, \text{ where } n \geq 0$$

$$= (11)^n 0, n \geq 0$$

I
 $|$
 $11I$
 $|$
 \circ

I
 $|$
 $11I$
 $|$
 $111I$
 $|$
 \circ

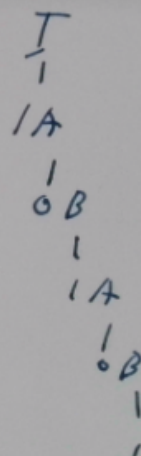
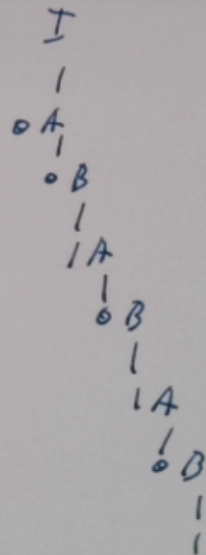
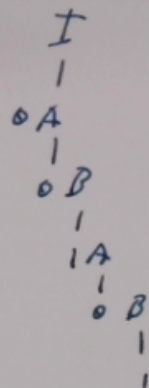
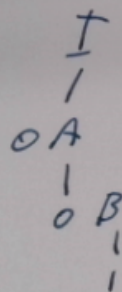
$$L(G) = \{1^n 0 : n \geq 0\}$$

$$V = \{0, 1, A, B, I\}, \quad T = \{0, 1\}, \quad S = I, \quad G = (V, T, S, P)$$

$$P = \{I \rightarrow 0A, I \rightarrow 1A, A \rightarrow 0B, B \rightarrow 1A, B \rightarrow 1\}$$

$$L(G) = \left\{ 001, 00101, 0010101, 10101, 101, 1010101, \dots \right\}$$

$$0(01)^n \text{ or } 1(01)^n \quad n \geq 1$$



Constructing a Grammar

$$L(G) = \{ \lambda, 01, 0011, 000111, 00001111, \dots \} = \begin{matrix} \cancel{(01)^n}, n \geq 0 \\ (0)^n (1)^n, n \geq 0 \end{matrix}$$

$$V = \{0, 1, I\} \quad T = \{0, 1\} \quad S = I$$

$$P = \{ I \rightarrow \lambda, \cancel{I \rightarrow 01}, I \rightarrow 0I1 \}$$

$$L(G) = 0^m 1^n, m, n \geq 0$$

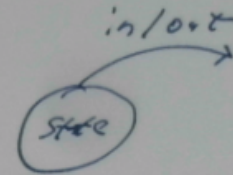
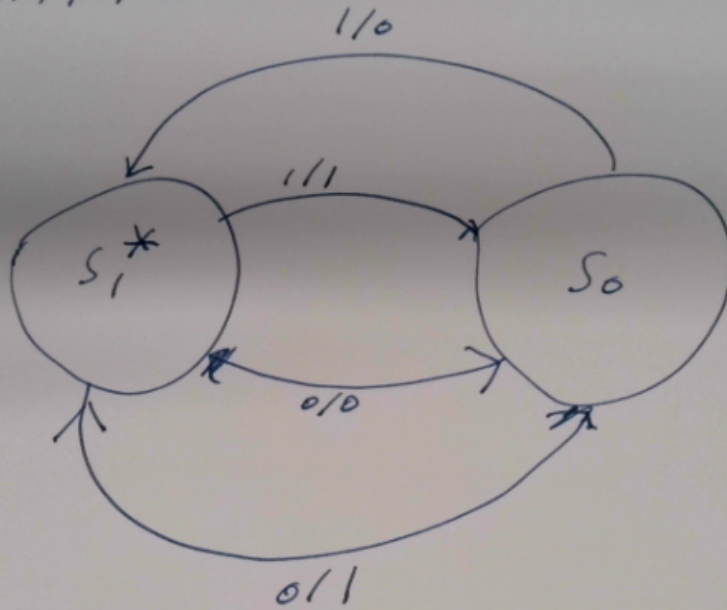
$$V = \{0, 1, I, A, B\}, T = \{0, 1\} \quad S = I$$

$$P = \{ A \rightarrow \lambda, A \rightarrow 0A, B \rightarrow \lambda, B \rightarrow 1B, I \rightarrow AB \}$$

Finite State Machines (FSM)

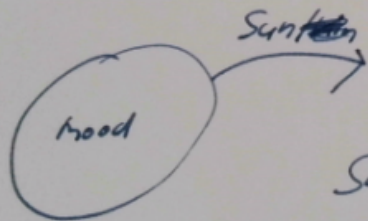
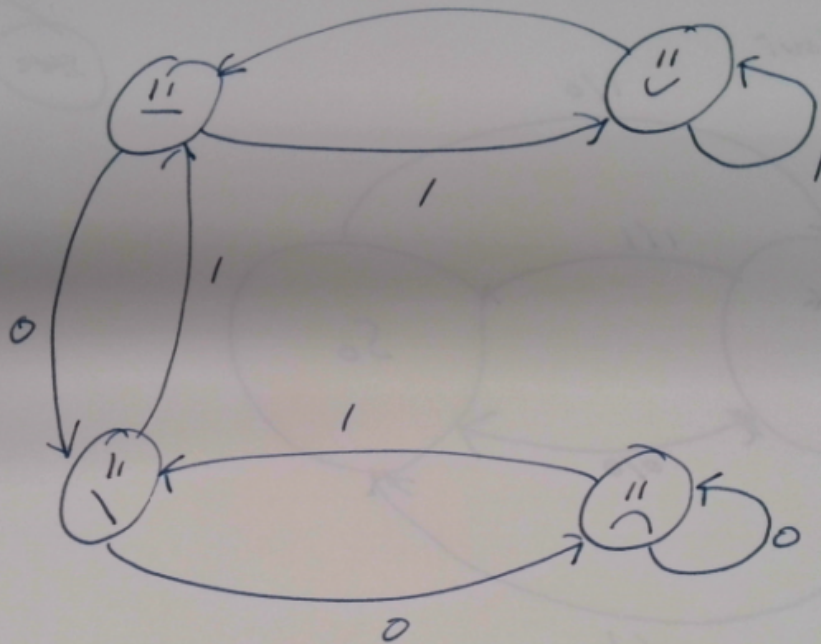
PS / IM / NS / out

Example:



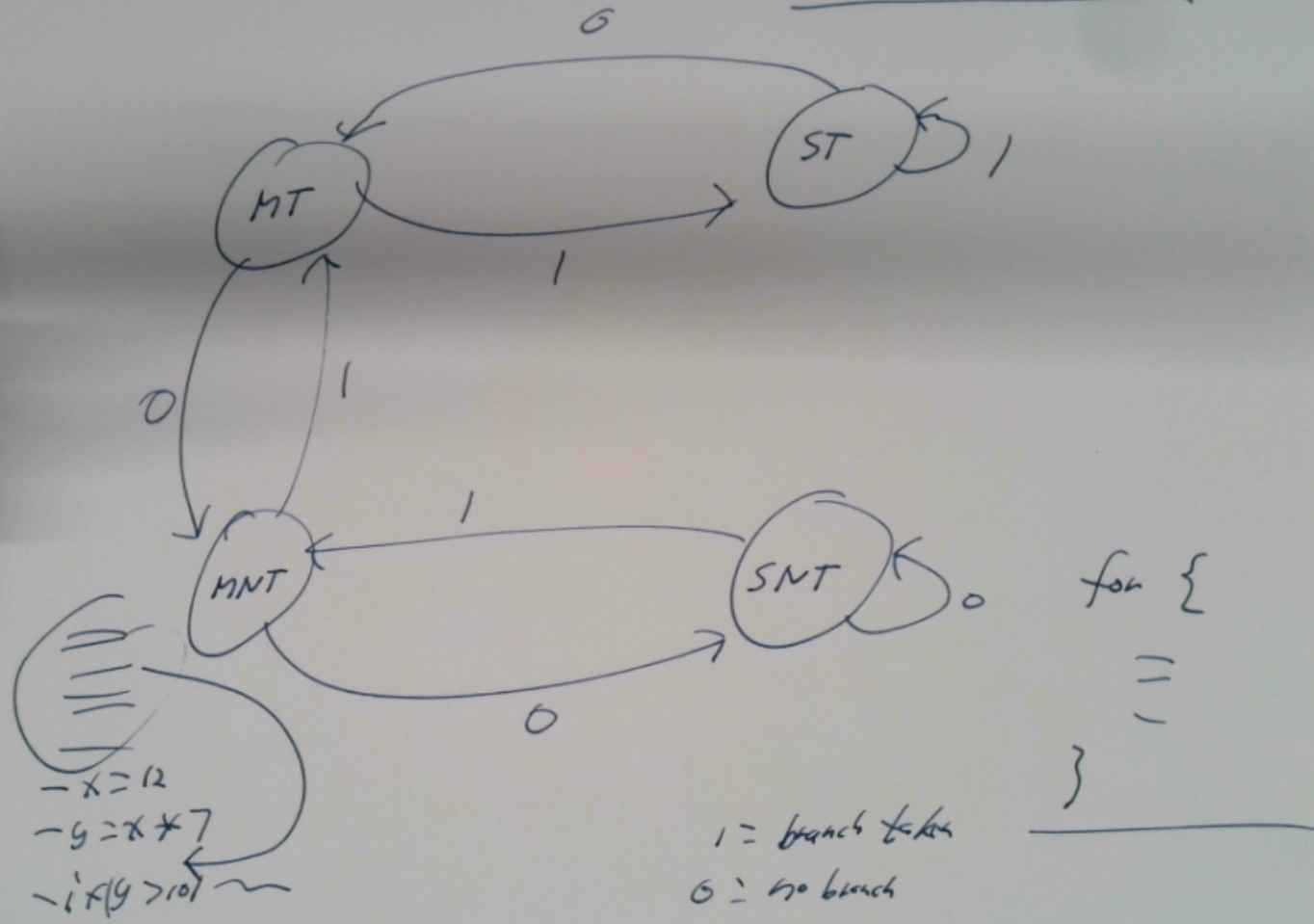
* initial state

in: 11001001
 out: 10011100
 ↑ ↑ ↑ ↑



Sun : 1 = Sah
 0 = Bah

BRANCH PREDICTOR



BRANCH PREDICTOR

